

Tutorial 2

Exercise 1: Consider the following atomic propositions:

- p — „the sun is shining“
- q — „it is raining“
- r — „we can see a rainbow“
- s — „it is snowing“

Describe in a natural language what the propositions represented by the following formulas of propositional logic say:

- a) $(p \wedge q) \rightarrow r$
- b) $p \rightarrow (\neg q \wedge \neg s)$
- c) $\neg\neg r$
- d) $(p \vee q) \vee s$
- e) $q \rightarrow q$
- f) $\neg s \rightarrow q$
- g) $(\neg r \wedge q) \leftrightarrow \neg s$
- h) $\neg(\neg p \rightarrow \neg q)$

Exercise 2: Represent the following propositions by formulas of propositional logic (for each formula, specify precisely what are the atomic propositions):

- a) If it is not Monday today then it won't be Wednesday after tomorrow.
- b) If is Monday or Wednesday today and it not Friday after tomorrow, then it is Monday today.
- c) It is not Monday nor Thursday today.

Determine the days in a week, on which these propositions are true, and on which are false.

Exercise 3: Write the following propositions as formulas of propositional logic (for each formula specify what are the atomic propositions):

- a) If barometric pressure drops, then it will be raining or snowing.
- b) If a packet with a request comes, this request will be processed and a packet with an acknowledgement will be sent, or a packet with information about error will be sent.
- c) If new oilfields are not found and there is a crisis in the Middle East, then oil prices will increase.
- d) If Mr. Smith has bought a new car and has not sold the old one, then he has already payed off his mortgage or he has got a new loan.
- e) Sister has a blue coat and a white coat.
- f) If John testifies and tells the truth, he will be found guilty; and if he does not testify, he will be found guilty.
- g) A sufficient condition for a number x to be odd is that x is a prime and it is greater than 2.
- h) Necessary condition for a sequence to converge is that it is bounded from above and from below.
- i) This amount will be paid if and only if the goods will be delivered.
- j) If x is positive, then x^2 is positive.

- k) If triangle ABC is not isosceles then it is not equilateral.
- l) Graph G is planar if and only if it does not contain as a subgraph a subdivision of graph K_5 nor a subdivision of graph $K_{3,3}$.
- m) It is not true that if this candidate won't be elected for president then the economical situation does not get worse.
- n) If the culprit forged this document, bribed the taxi-driver, and haven't cleared the fingerprints, then an evidence will be found against him.

Exercise 4: Consider the following propositions:

- p — „Prague is larger than Liberec“
 q — „Carlsbad is situated in western Bohemia“
 r — „the Elbe flows through České Budějovice“

(So propositions p and q are true, and proposition r is false.)

Which of the following propositions are true, and which are false? (Formulate these propositions also in a natural language.)

- | | |
|--|--|
| a) $p \vee r$ | e) $(q \vee \neg r) \rightarrow p$ |
| b) $p \wedge r$ | f) $(q \vee p) \rightarrow (q \rightarrow \neg r)$ |
| c) $\neg p \wedge \neg r$ | g) $(q \leftrightarrow \neg p) \leftrightarrow (p \leftrightarrow r)$ |
| d) $p \leftrightarrow (\neg q \vee r)$ | h) $(q \rightarrow p) \rightarrow ((p \rightarrow \neg r) \rightarrow (\neg r \rightarrow q))$ |

Exercise 5: For each of the following sequences of symbols, do the following:

- a) Decide if it is a well-formed formula of propositional logic (according to the formal definition).
- b) Decide if it is a well-formed formula of propositional logic when the conventions for omitting parentheses can be used.
- c) If it is a well-formed formula (either according to (a) or (b)):
- Write this formula according to formal definition (i.e., without omitting parentheses).
 - Write this formula with as much parentheses omitted as possible.
 - Draw a corresponding abstract syntax tree.

(Justify your answers in points (a) and (b).)

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|--------------------------------------|--|
| 1. $(p) \neg \wedge \wedge$ | 8. $\wedge pq$ |
| 2. $\forall x : q(x) \wedge r(x, x)$ | 9. $p \wedge q$ |
| 3. p | 10. $(p \wedge q)$ |
| 4. $(\neg(\neg q))$ | 11. $((p \wedge q))$ |
| 5. $(\neg(\neg q()))$ | 12. $((p \wedge q) \vee r)$ |
| 6. $(\neg(\neg)q)$ | 13. $((\neg p) \vee (q \leftrightarrow (\neg r)))$ |
| 7. $(p \neg q)$ | 14. $r \vee (\neg q \vee s)$ |

15. $((\neg r \vee \neg p) \vee s) \wedge (\neg q \vee s)$

16. $(\neg((\neg p) \rightarrow (\neg(\neg r))))$

Exercise 6: Using the table method, determine all models of the following formulas and decide, which of these formulas are tautologies, which are satisfiable, and which are contradictions:

a) $p \vee q$

b) $p \vee \neg p$

c) $p \vee q \rightarrow q \vee p$

d) $p \rightarrow (p \vee q) \vee r$

e) $p \rightarrow (\neg p \rightarrow q)$

f) $(p \rightarrow q) \rightarrow (q \rightarrow p)$

g) $((p \rightarrow q) \leftrightarrow q) \rightarrow p$

h) $p \rightarrow (q \rightarrow (q \rightarrow p))$

i) $p \wedge \neg(q \rightarrow p)$

j) $p \wedge q \rightarrow p \vee r$

k) $(p \vee (\neg p \wedge q)) \vee (\neg p \wedge \neg q)$

l) $p \wedge q \rightarrow (p \leftrightarrow q \vee r)$

m) $(p \wedge q \rightarrow (p \wedge \neg p \rightarrow q \vee \neg q)) \wedge (q \rightarrow q)$

n) $p \leftrightarrow q$

o) $p \leftrightarrow p \vee p$

p) $p \vee q \leftrightarrow q \vee p$

q) $(p \rightarrow q) \leftrightarrow (q \rightarrow p)$

r) $(p \leftrightarrow p) \leftrightarrow p$